

OWL and Description Logics DL

By:

A.Aziz Altowayan @ Pace University

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Description Logics (DL)

What:

- Knowledge **representation** language.
(serving primarily for formal description of concepts and roles (relations)).

Why:

- Used for formal **reasoning** on the concepts of a domain.

OWL's relation to DL:

"In the beginning, IS-A was quite simple. Today, however, there are almost as many meanings for this inheritance link as there are knowledge-representation systems."

Ronald Brachman 1983

DL Terminologies

In OWL:

- A **class** is a collection of objects.
- A **property** is a directed binary relation.
- An **instance** is an object.

In DL, the above corresponds to:

- A **concept**.
- A **role**.
- An **individual**.

A **concept** corresponds to a *unary predicate* while a **role** corresponds to a *binary predicate*.

Concept (Formulae):	e.g. Human, Male, Female, Animal
Roles (Modalities):	e.g. hasChild, hasParent, loves
Individuals (Ground term):	e.g. Aziz, Lixin, USA

DL modeling and Knowledge Base

KR based on DLs, consists of 2 components:

- **TBox**, "*Terminological Box*" (describes terminology).
 - contains sentences describing concept hierarchies (i.e., relations between concepts).
 - e.g.) Every employee is a person.
- **ABox**, "*Assertion Box*" (assertions about individuals).
 - contains ground sentences stating where in the hierarchy individuals belong (i.e., relations between individuals and concepts).
 - e.g.) Bob is an employee.

“Reasoning in ontologies and knowledge bases is one of the reasons why a specification needs to be formal one.”

Formal description and notations

Symbol	Description	Example	Read
\top	all concept names	\top	top
\perp	<u>empty</u> concept	\perp	bottom
\sqcap	<u>intersection</u> or <u>conjunction</u> of concepts	$C \sqcap D$	C and D
\sqcup	<u>union</u> or <u>disjunction</u> of concepts	$C \sqcup D$	C or D
\neg	<u>negation</u> or <u>complement</u> of concepts	$\neg C$	not C
\forall	<u>universal restriction</u>	$\forall R.C$	all R-successors are in C
\exists	<u>existential restriction</u>	$\exists R.C$	an R-successor exists in C
\sqsubseteq	Concept <i>inclusion</i>	$C \sqsubseteq D$	all C are D
\equiv	Concept <i>equivalence</i>	$C \equiv D$	C is equivalent to D
\doteq	Concept <i>definition</i>	$C \doteq D$	C is defined to be equal to D
:	Concept <i>assertion</i>	$a : C$	a is a C
:	Role <i>assertion</i>	$(a, b) : R$	a is R-related to b

See:

- [DL wiki](#)

OWL syntax, DL syntax, and semantics

OWL abstract syntax	DL syntax	Semantics
<i>Class descriptions</i>		
Class (A)	A	$A^{\mathcal{I}} \subseteq \Delta^{\mathcal{I}}$
owl:Thing	\top	$\top^{\mathcal{I}} = \Delta^{\mathcal{I}}$
owl:Nothing	\perp	$\perp^{\mathcal{I}} = \emptyset$
intersectionOf ($C_1 \dots C_n$)	$C_1 \sqcap \dots \sqcap C_n$	$(C_1 \sqcap \dots \sqcap C_n)^{\mathcal{I}} = C_1^{\mathcal{I}} \cap \dots \cap C_n^{\mathcal{I}}$
unionOf ($C_1 \dots C_n$)	$C_1 \sqcup \dots \sqcup C_n$	$(C_1 \sqcup \dots \sqcup C_n)^{\mathcal{I}} = C_1^{\mathcal{I}} \cup \dots \cup C_n^{\mathcal{I}}$
complementOf (C)	$\neg C$	$(\neg C)^{\mathcal{I}} = \Delta^{\mathcal{I}} \setminus C^{\mathcal{I}}$
oneOf($w_1 \dots w_n$)	$\{w_1, \dots, w_n\}$	$(w_1, \dots, w_n)^{\mathcal{I}} = \{w_1^{\mathcal{I}}, \dots, w_n^{\mathcal{I}}\}$
restriction (P someValuesFrom(E))	$\exists P.E$	$(\exists P.E)^{\mathcal{I}} = \{a \mid \exists w. \langle a, w \rangle \in P^{\mathcal{I}} \wedge w \in E^{\mathcal{I}}\}$
restriction (P allValuesFrom(E))	$\forall P.E$	$(\forall P.E)^{\mathcal{I}} = \{a \mid \forall w. \langle a, w \rangle \in P^{\mathcal{I}} \rightarrow w \in E^{\mathcal{I}}\}$
restriction (P hasValue(w))	$\exists P.\{w\}$	$(\exists P.w)^{\mathcal{I}} = \{a \mid \langle a, w^{\mathcal{I}} \rangle \in P^{\mathcal{I}}\}$
restriction (P minCardinality(n))	$\geq nP$	$(\geq nP)^{\mathcal{I}} = \{a \mid \#\{w \mid \langle a, w \rangle \in P^{\mathcal{I}}\} \geq n\}$
restriction (P maxCardinality(n))	$\leq nP$	$(\leq nP)^{\mathcal{I}} = \{a \mid \#\{w \mid \langle a, w \rangle \in P^{\mathcal{I}}\} \leq n\}$
restriction (P cardinality(n))	$= nP$	$(= nP)^{\mathcal{I}} = (\geq nP \cap \leq nP)^{\mathcal{I}}$
$P \in \{R, T\}$ $w \in \{a, v\}$ $E \in \{C, d\}$		

See:

- complete list
- [DL syntax and semantic](#)

DL ALC and its concepts

Attributive concept Language with Complements ALC is a member of DL family, where:

- **top** is a concept.
- **bottom** is a concept.
- all **atomic concepts** are concepts
- the **intersection** of two concepts is a concept
- the **union** of two concepts is a concept
- the **complement** of a concept is a concept
- the **universal restriction** of a concept by a role is a concept
- the **existential restriction** of a concept by a role is a concept

Inference in DL (Decision problems)

“Description logics are created with the focus on tractable reasoning.”

Examples of tasks required from reasoners are:

- Instance checking. *(is a an instance of A?)*
- Relation checking. *(is a related to b?)*
- Subsumption. *(is A a subset of C?)*
- Concept consistency. *(is there any contradiction in A's definition?)*

Finally, Tradeoff between Expressive Power and Efficient Reasoning Support.

The **richer** the language is, the **more inefficient** the reasoning support becomes.